

Design of A Low Power Area Optimized 4-Bit Arithmetic Logic Unit for High Speed Processors

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Abstract: With the advancement in technology area, power and speed had become a major concern for the growing VLSI industry. Static CMOS logic does not fully meet the needs of future computing. Therefore improved digital logic techniques and styles which are energy efficient, fast and consume lesser area must be utilized. In this paper, a low power, area optimized and high speed 4-Bit Arithmetic Logic Unit having minimum circuit complexity is presented in an elegant way. The major components of the ALU are designed using Modified Gate Diffusion Input (MGDI) technique. The proposed ALU is designed using Cadence Virtuoso tool in 90nm process technology. The simulations are conducted using the Spectre simulation tool at a supply voltage of 1.8V. The simulation results show that the proposed design has lower delay, consumes less power and requires less area.

Keywords: GDI, MGDI, ALU, Power, Delay.

I. Introduction

In the era of growing technology and scaling of devices up to nanometer regime, the demand for low power and high speed microelectronic devices has come to the forefront. Nowadays there is a rapid increase in the utilization of portable applications which demands small-size, low power, high speed and high throughput circuitry. One of the most important parts of a digital processor is the arithmetic logic unit. ALU is designed to do the arithmetic and logic operations which are the basic processes that need to be done for almost any data that is being processed by central processing unit (CPU). And that's why the ALU is called heart of microprocessor, microcontroller and digital signal processor. The performance of high speed processing system completely depends on the performance of the ALU and hence it is required that for faster calculations the ALU must be fast. In this paper, a 4-Bit ALU has been designed using a low power 8T hybrid full adder. The sub-circuits i.e. multiplexers, OR Gate, AND gate and inverter are designed using the MGDI technique which requires lesser transistor count.

The paper is organized as follows: Section II gives an overview of previous works done. Section III gives a brief account of Gate Diffusion Input (GDI) technique and Modified Gate Diffusion Input (MGDI) technique. In section IV the components that are required to construct the ALU and its operations are explained. Section V shows the simulation results and performance analysis of the proposed ALU. Section VI concludes the paper.

II. Previous Works

Static CMOS logic is the most commonly used logic design technique. In this, circuits are made up of two networks namely pull up network (PUN) and pull down network (PDN). The main drawback of static CMOS logic is that it uses more number of PMOS transistors because of which the delay and area of the circuit is increased. In some scenarios, the driving capability of the circuit reduces [1][2]. A 1-Bit ALU is implemented using different foundries like 45nm, 65nm and 90 nm. The performance of developed ALU has been analyzed and compared in terms of area and power using BSIM4 device model [3]. A high-speed 4-bit ALU has been designed for 1V operation to demonstrate the usefulness of the back-gate forward substrate bias (BGFSB) method in 1.2 μ m n-well CMOS technology [4]. A comparison between the number of transistors used in traditional CMOS technique and GDI design of ALU is done and an inference is drawn that GDI design evidently reduces the number of transistors and hence optimizes the area of ALU as well increase its working speed [5]. An area optimized low power hybrid Arithmetic Logic Unit is implemented using pass transistor logic and GDI technique [6]. When compared to traditional CMOS technique GDI is easier to implement and efficient in terms of transistor count [7].

III. Overview of GDI And MGDI Techniques

This section gives a brief description of Gate Diffusion Input (GDI) technique and Modified Gate Diffusion Input (MGDI) technique.

3.1. Gate Diffusion Input (GDI) Technique

The structure of GDI cell was first proposed by Morgenshtein [8]. GDI is a new technique of realizing CMOS static logic technique, which is based on the use of a simple cell structure as shown in fig. 1. This technique is used in digital combinational circuit cell design because of its low power consumption and minimum area on the chip. At first sight a GDI cell looks similar to that of a standard CMOS inverter but there are only few differences. In the place of VDD and GND pins, input signals are used which makes it more flexible than CMOS design. The GDI cell has three inputs namely: G- common input to the gate of NMOS and PMOS; N- input to the source/drain of NMOS; P- input to the source/drain of PMOS. Bulks of both PMOS and NMOS are connected to P or N respectively, so it will be self biased at contrast to CMOS inverter

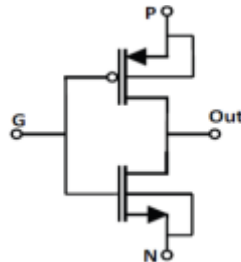


Fig. 1: Basic GDI Cell

3.2. Modified Gate Diffusion Input Technique

The basic MGDI cell as shown in fig. 2 is similar to that of GDI cell which consists of NMOS and PMOS containing four terminals: 'G' is the common gate input of NMOS and PMOS transistors, 'P' is the outer diffusion node of PMOS transistor, 'N' is the outer diffusion node of MOS transistor, and 'Out' is the common diffusion node of both transistors.

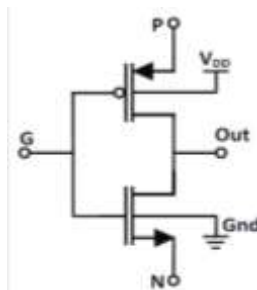


Fig. 2: Basic MGDI Cell

In the MGDI cell, the bulk node of PMOS transistor is connected to the supply voltage VDD and the bulk of NMOS transistor is connected to ground (GND). This makes the GDI to constant body biasing in modified GDI Cell which in turn increases the stability of the circuit and its loading effect. The MGDI cell uses standard four-terminal NMOS and PMOS transistors and can be effortlessly implemented in all types of standard CMOS technology.

Table 1: Basic Functions Using GDI/MGDI Cell

N	P	G	OUT	FUNCTION
0	1	A	A'	Inverter
0	B	A	A'B	Function1
B	1	A	A'+B	Function2
1	B	A	A+B	OR
B	0	A	AB	AND
C	B	A	A'B+AC	Multiplexer

Table 1 shows different Boolean operations realized just by changing the control inputs to the cell. A large number of functions can be implemented using the basic GDI/MGDI cell. Function1 (F1) and Function2 (F2) are the universal sets for GDI/MGDI which can be implemented by using only 2 transistors.

IV. Proposed Work

ALU is one of the most important parts of a digital computer which is designed to do the arithmetic and logic operations that need to be done for almost any data that is being processed by a processor. This section covers the building blocks and the operation of the proposed Design.

4.1. Multiplexer

Multiplexer is a digital switch which has number of input data lines and a single output line. Selection line plays a major role to select a particular input. There are '2ⁿ' input lines and 'n' selection lines whose bit combinations determine which input is to be transferred to the output. Fig. 3 and fig. 4 represents the schematics of 2x1 multiplexer and 4x1 multiplexer respectively.

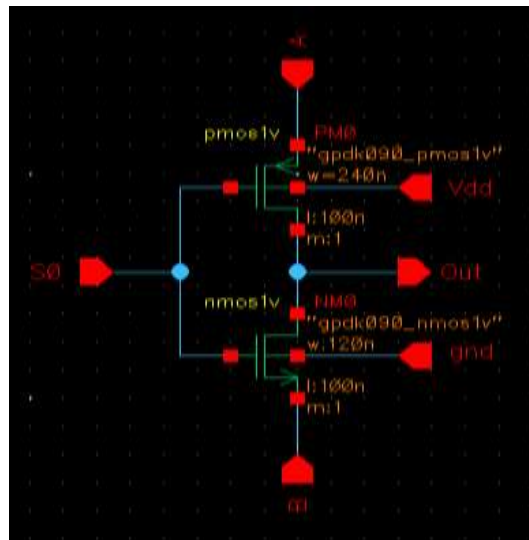


Fig. 3: Schematic of 2x1 Multiplexer

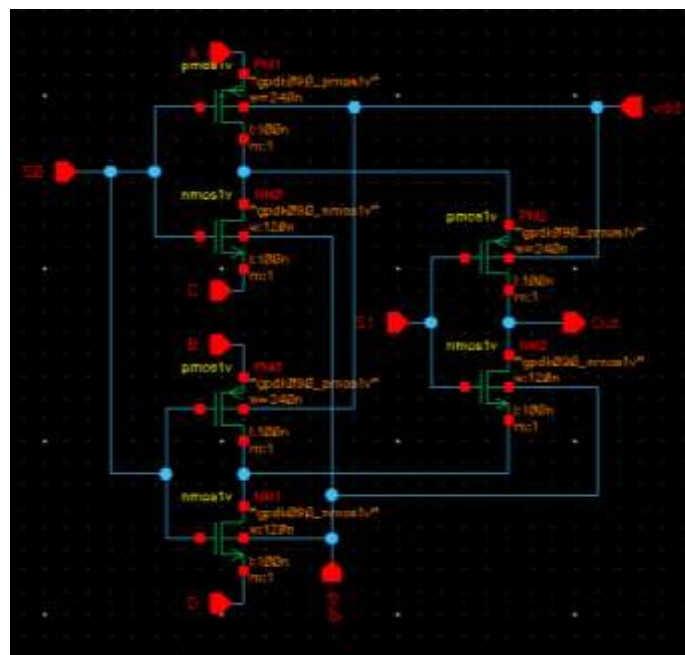


Fig. 4: Schematic of 4x1 Multiplexer

4.2. Logic Gates

Logic gates are the building blocks of an ALU which are responsible for all the logical operations in the circuit. The logic block in the proposed circuit consists of AND gate, OR gate, EXNOR gate and an inverter. The schematics of an inverter, AND gate, OR gate and EXNOR gate are shown in fig. 5, fig. 6, fig. 7 and fig. 8 respectively.

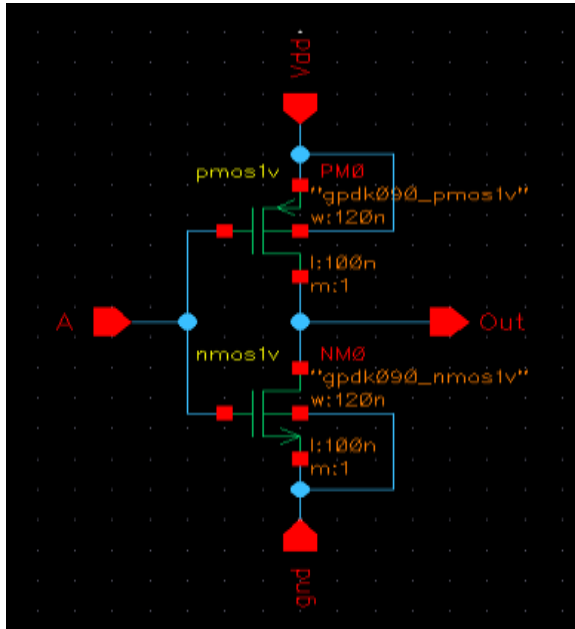


Fig. 5: Schematic of an Inverter

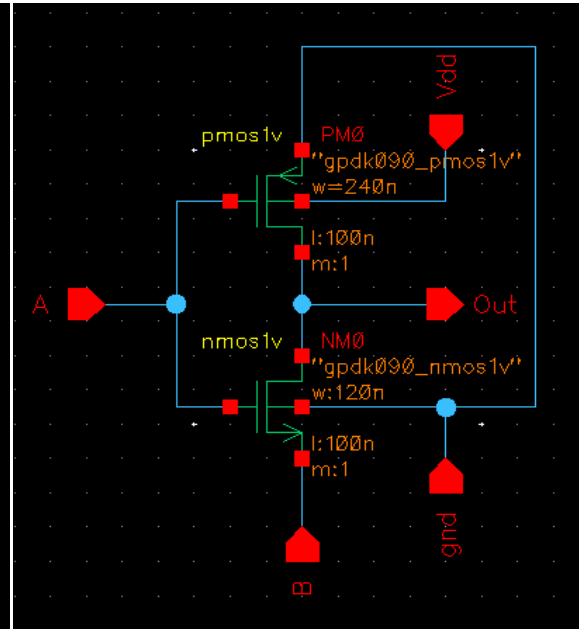


Fig. 6: Schematic of AND Gate

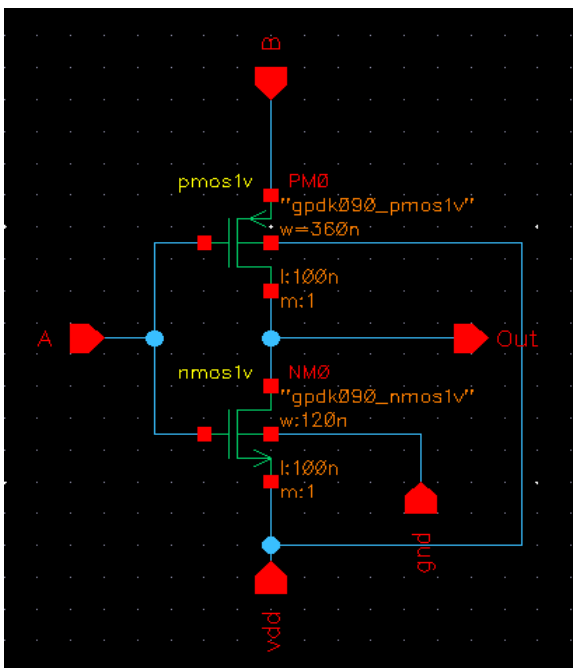


Fig. 7: Schematic of OR Gate

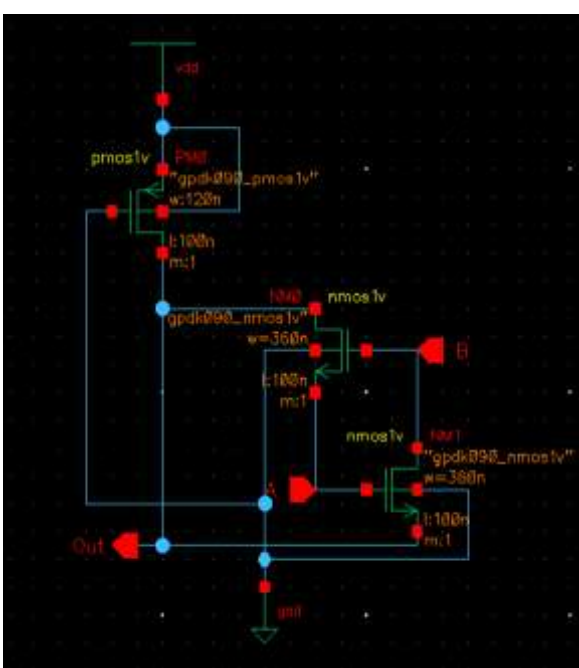


Fig. 8: Schematic of EXNOR Gate

3.4. Full Adder

Full adders are the basic elements of a large number of complex arithmetic circuits like multipliers, dividers, ALUs, etc. Therefore, improvement in the design of full adder circuit results in the overall performance boosting of the entire system. Schematic of the Full Adder (8T) in the proposed design is shown in fig. 9.

A full adder circuit adds binary numbers and accounts for values carried in. A single bit full adder operates on single-bit numbers and adds them. This is generally written as A, B, and C_{in} . A and B are operands of the addition operation whereas C_{in} bit is carried out in from a consequent less significant stage. The circuit produces a double-bit output i.e. Sum and Carryout (C_{out})

where,

$$\text{Sum} = A \oplus B \oplus C_{in} \quad (1)$$

$$C_{out} = AB + AC_{in} + BC_{in} \quad (2)$$

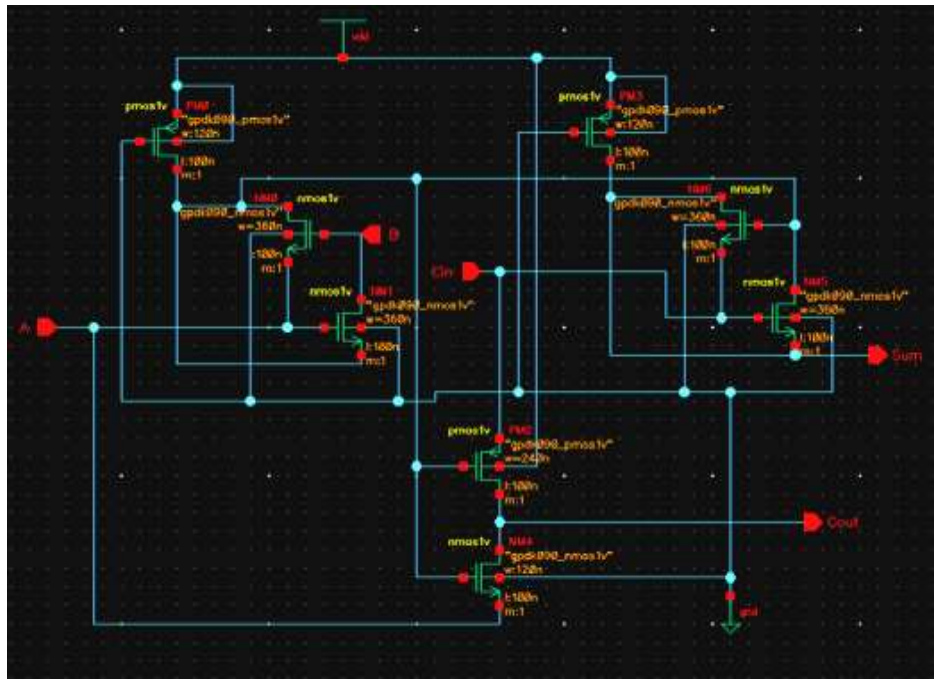


Fig. 9: Schematic of 1-Bit Full Adder (8T)

4.4. Design of Proposed 4-bit ALU

The proposed 4-bit ALU design consists of eight 4x1 multiplexers, four 8T full adders, four logic blocks and four 2x1 multiplexers. The output carry C_{i+1} of an earlier stage is feed to the input carry C_i of the next stage. The block diagram of the proposed 4-bit ALU is shown in fig. 10.

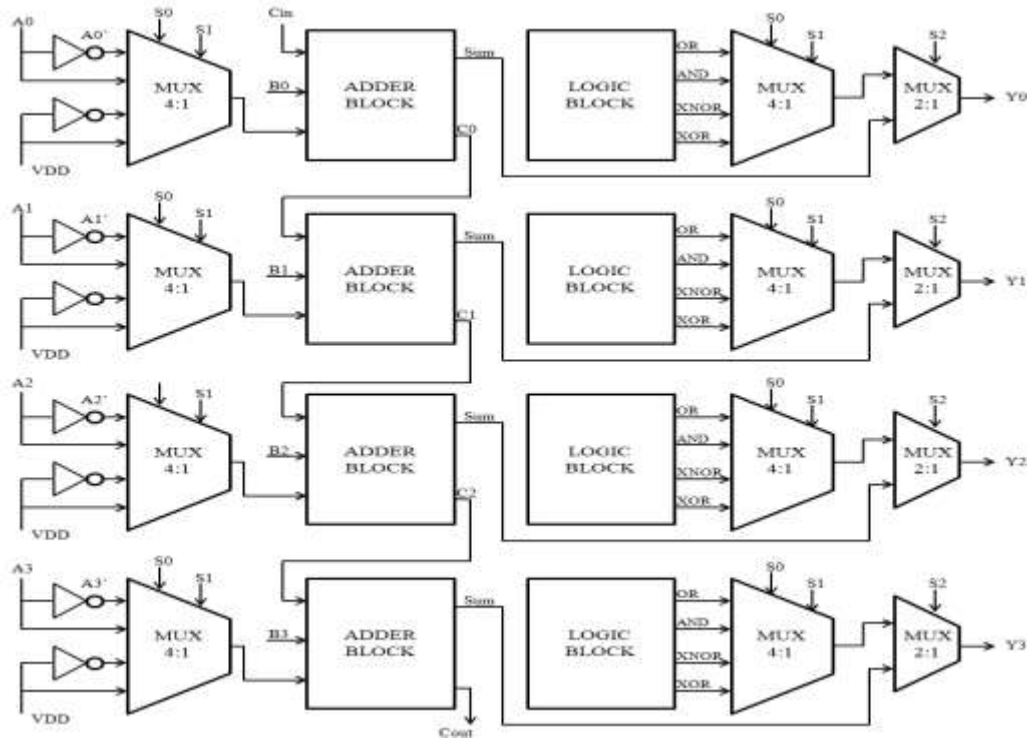


Fig. 10: Block Diagram of 4-Bit ALU

4.5. Operation of Proposed 4-bit ALU

The twelve micro-operations that are performed by the ALU are listed in table 2. The logic operations are selected with the selection line $S_2 = 0$. The input carry C_{in} has no effect during the logic operations and is therefore marked with don't care. The arithmetic operations are selected with the selection line $S_2 = 1$.

Table 2: Truth Table for 4-Bit ALU

S_2	S_1	S_0	C_{in}	Micro-Operation	Function
0	0	0	x	$A \vee B$	OR
0	0	1	x	$A \oplus B$	XNOR
0	1	0	x	$A \wedge B$	AND
0	1	1	x	$A \oplus B$	XOR
1	0	0	0	$A' + B$	Subtraction with borrow
1	0	0	1	$A' + B + 1$	Subtraction
1	0	1	0	B	Transfer B
1	0	1	1	$B + 1$	Increment B
1	1	0	0	$A + B$	Addition
1	1	0	1	$A + B + 1$	Addition with carry
1	1	1	0	$B - 1$	Decrement B
1	1	1	1	B	Transfer B

The arithmetic micro-operations are realized by controlling the data input to the full adder. For the Subtraction operation complement of A is used. The full adder performs this subtraction operation by two's complement method. For the Increment and Decrement operations, 'logic 1' and 'logic 0' are used as inputs respectively.

Logic micro-operations are performed by the logic block. The outputs of the gates are applied to the data inputs of the 4x1 multiplexer. The selection inputs S_1 and S_0 opt one of the data inputs of the multiplexer and direct its values to the output.

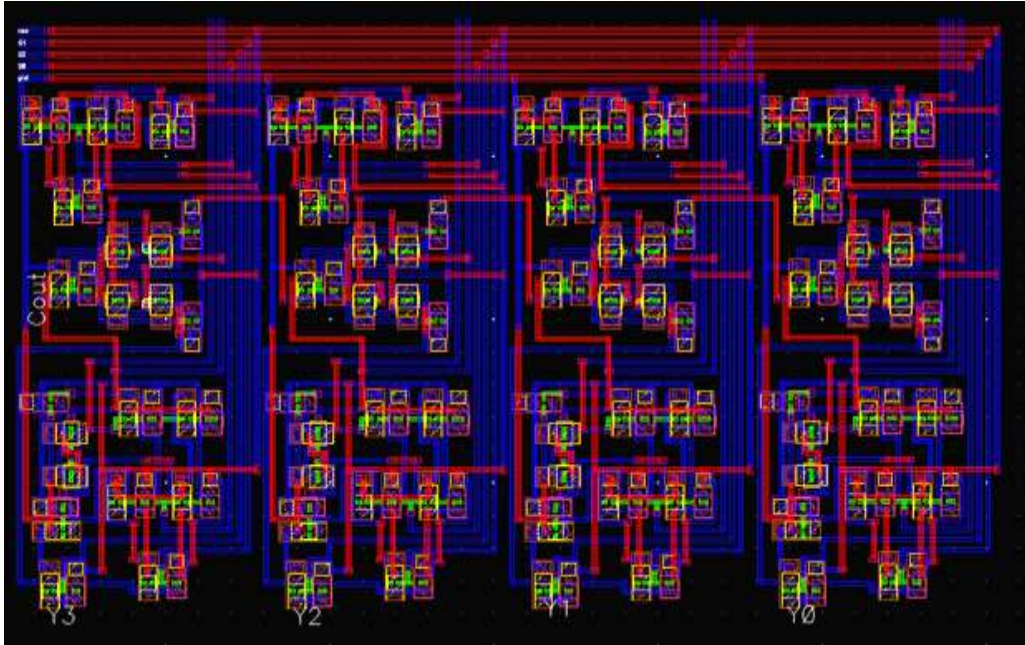


Fig. 11: Layout of Proposed 4-BIT ALU

The layout of the proposed 4-bit ALU is shown in fig. 11. The layout has been verified by performing all the requisite checks like the DRC checks (Design Rule Check), ERC checks (Extracted Resistance and Capacitance) and LVS checks (Layout versus Schematic) and is free from errors.

V. Results and Analysis

The proposed 4-Bit ALU is designed using Cadence Virtuoso tool in 90nm process technology. To obtain best power and delay performance the size of the transistors is adjusted without compromising with the output voltage swing. All the pre-layout and post-layout simulations are done using the Spectre based Cadence Virtuoso simulator with a power supply of 1.8V.



Fig. 12: Post Layout Simulation of 4-Bit ALU

The post layout simulation of the proposed 4-Bit ALU is shown in fig. 12. The simulation is performed in Virtuoso Analog Design Environment (ADE) using Spectre tool with the stop time of 240 ns. For simplification, the input signals $A = [A_0 A_1 A_2 A_3]$ and $B = [B_0 B_1 B_2 B_3]$ are of constant value throughout the period i.e., $A_0 = A_1 = A_2 = A_3$ and $B_0 = B_1 = B_2 = B_3$. The ALU performs different operations with the alteration in its control. The overall performance parameters of the proposed 4-Bit ALU using the 90 nm process technology at a supply voltage of 1.8V are shown in table 3.

Table 3: Performance Parameters of Proposed 4-Bit ALU

Transistor Count	Average Power (W)	Delay (s)	Area (m ²)
132	913×10^{-6}	233.6×10^{-12}	540.74×10^{-12}

VI. Conclusion

A low power area optimized 4-Bit ALU designed in 90 nm process technology using the MGDI technique and simulated using the Cadence Virtuoso based Spectre simulator. The hybrid Full adder constructed with only eight transistors utilizing XNOR-XNOR logic and a multiplexer has a better output voltage swing. The proposed design comprises of only 132 transistors and operates at 1.8V supply voltage and has a power consumption of $913 \mu\text{W}$. The area on chip is $540.74 \times 10^{-12} \text{ m}^2$ and has a delay of 233.6 ps. Overall simulation results shows that the proposed 4-Bit ALU has minimum delay, consumes less power and has less transistor count when compared to other design styles.

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